## L7s Multiple Output example

## The Problem (from text Unit 14)

## - Problem Statement

14.5 A sequential circuit has one input $(X)$ and two outputs $\left(Z_{1}\right.$ and $\left.Z_{2}\right)$. An output $Z_{1}$ occurs every time the input sequence 010 is completed, provided that the sequence in has never occurred. An output $Z_{2}=1$ occurs every time the input 100 is completeil Note that once a $Z_{2}=1$ output has occurred, $Z_{1}=1$ can never occur but not vice venii Find a Mealy state graph and state table (minimum number of states is eight).

- Will do for Mealy and Moore
- One input X, two outputs Z1 and Z2
- Z1 = 1 occurs every time 010 is last 3 on input, provided 100 has never occurred
- Z2 = 1 every time 100 is last 3 on input


## Choose a starting state (Mealy)

$\square$ This is the state after a reset.
$\square$ The slides will show the progression (developed on the board - now slides)

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| State | Meaning |
| :--- | :--- |
| S0 | Starting State |

}

## Now add states transition from S0

$\square$ When in S0 what happens when a 1 is input or a 0 is input?

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
|  |  |
| S3 | an initial 1 of possibly 100 |



## In S1

$\square$ What happens on input of 0 - stay in S1
$\square$ What happens on input of 1 - transition to new state S2.

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
| S2 | have 01 as last 2 inputs |
| S3 | an initial 1 of possibly 100 |



## In S3

## - Input of 1 - stay in S3

- Input of 0 - now have 10 of possible 100 seq

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
| S2 | have 01 as last 2 inputs |
| S3 | an initial 1 of possibly 100 |
| S4 | have 10 as last 2 inputs |



## In S2

- Input 1 - have 11 as last two $-1^{\text {st }} 1$ of 100 - transition to S3
- Input 0 - have completed 010 and have 10 as last two inputs, S4

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
| S2 | have 01 as last 2 inputs |
| S3 | an initial 1 of possibly 100 |
| S4 | have 10 as last 2 inputs |



## In S4

- Input 1 - have 01 as last 2 inputs - go to S2
- Input 0 - have 100 as last 3 and sequence 100 - go to new state S 5 and sequence 010 can not be recognized again

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
| S2 | have 01 as last 2 inputs |
| S3 | an initial 1 of possibly 100 |
| S4 | have 10 as last 2 inputs |
| S5 | have seen sequence 100 |



## In S5

- Input 0 - Have 00 as last two - not start of 100
- Input 1 - Have possible start of 100 - transition to S6

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
| S2 | have 01 as last 2 inputs |
| S3 | an initial 1 of possibly 100 |
| S4 | have 10 as last 2 inputs |
| S5 | have seen sequence 100 |
| S6 | have a $1-$ start of 100 |



## In S6

## - Input 0 - have 10 of 100 - go to S7

- Input 1 - Stay in S6

State Meaning<br>Starting State<br>an initial 0 of possibly 010<br>have 01 as last 2 inputs an initial 1 of possibly 100 have 10 as last 2 inputs have seen sequence 100 have a 1 - start of 100 have a 10 of possible 100



## In S7

- Input 0 - Now have 100 and output a 1 - go to S5
- Input 1 - Have a $1^{\text {st }} 1$ and could be start of 100 - go to S6
- Done

| State | Meaning |
| :--- | :--- |
| S0 | Starting State |
| S1 | an initial 0 of possibly 010 |
| S2 | have 01 as last 2 inputs |
| S3 | an initial 1 of possibly 100 |
| S4 | have 10 as last 2 inputs |
| S5 | have seen sequence 100 |
| S6 | have a 1 - start of 100 |
| S7 | have a 10 of possible 100 |



## Moore Machine implementation

- The Moore Machine implementation
- It adds 2 more states and is left to the student to work this through. Remember that the output is associated with the state, not a combination of the state and input as in a Mealy Machine.


## State S0

$\square$ Have a starting state S0 and its meaning
$\square$ Remember - a Moore machine

State Meaning<br>S0 Starting state

ㅁ Note output designation on State symbol

## In S0

## - 0 input - have the start of 010 <br> $\square 1$ input - have the start of 100

State Meaning<br>S0 Starting state<br>S1 have $1^{\text {st }} 0$ of start of 010<br>S3 have $1^{\text {st }} 1$ of start of 100



## In S1

- 0 input - last 2 are 00 - stay in S1
$\square 1$ input - last 2 are 01 - transition to state S2
State Meaning
S0 Starting state
S1 have $1^{\text {st }} 0$ of start of 010
S2 have 01 as last 2 inputs
S3 have $1^{\text {st }} 1$ of start of 100
S4 have 10 as last 2 inputs
S5 010 detected -10 as last two inputs
S6 100 detected - output Z2 $=1$
S7 after 100 - a 0 input
S8 after 100 - a 1 input
S9 after 100 - have 10 as last 2


## In S3

## - 0 input - have 10 as last 2 - go to S4

## s: hemput - have 11 as last 2 - stay in S3

S0 Starting state
S1 have $1^{\text {st }} 0$ of start of 010
S2 have 01 as last 2 inputs
S3 have $1^{\text {st }} 1$ of start of 100
S4 have 10 as last 2 inputs
S5 010 detected - 10 as last two inputs
S6 100 detected - output $\mathrm{Z} 2=1$
S7 after 100 - a 0 input
S8 after 100 - a 1 input


S9 after 100 - have 10 as last 2

## In S2

## - 0 input - Have 010 as last 3 - 10 as last 2 - go to state S 5 which has $\mathrm{Z} 1=1$ as its output

$\square_{\frac{\text { state }}{\text { so }}} \underset{\substack{\text { intang } \\ \text { Stating state }}}{ }$ go to S3 as 11 are last 2 inputs
S1 have $1^{\text {st }} 0$ of start of 010
S2 have 01 as last 2 inputs
S3 have $1^{\text {st }} 1$ of start of 100
S4 have 10 as last 2 inputs
S5 010 detected - 10 as last two inputs
S6 100 detected - output $\mathrm{Z} 2=1$
S7 after 100 - a 0 input
S8 after 100 - a 1 input
S9 after 100 - have 10 as last 2


## In S4

## - 0 input - 100 has been detected - new state S6 where 010 can not be detected output Z2=1

- 1 input - last 3 are 101, i.e., last 2 are 01 - go sato $\mathrm{ta}_{1}$ S2
S0 Starting state
S1 have $1^{\text {st }} 0$ of start of 010
S2 have 01 as last 2 inputs
S3 have $1^{\text {st }} 1$ of start of 100
S4 have 10 as last 2 inputs
S5 010 detected - 10 as last two inputs
S6 100 detected - output Z2 $=1$
S7 after 100 - a 0 input
S8 after 100 - a 1 input
S9 after 100 - have 10 as last 2



## In S5 - 010 detected

## - 0 input - 100 are last 3 - go to S6 <br> - 1 input - 101 are last 3, 01 last 2 - go to S2

| State | Meaning |
| :--- | :--- |
| S0 | Starting state |
| S1 | have $1^{\text {st }} 0$ of start of 010 |
| S2 | have 01 as last 2 inputs |
| S3 | have $1^{\text {st }} 1$ of start of 100 |
| S4 | have 10 as last 2 inputs |
| S5 | 010 detected -10 as last two inputs |
| S6 | 100 detected - output Z2 = 1 |
| S7 | after $100-$ a 0 input |
| S8 | after $100-$ a 1 input |
| S9 | after $100-$ have 10 as last 2 |



## In S6

## - S6 have detected 100 and output Z2 = 1 <br> - 0 input - new state S 7 - means a 0 input <br> ㅁ 1 input - new state S 8 - means a 1 received

State Meaning<br>S0 Starting state<br>S1 have $1^{\text {st }} 0$ of start of 010<br>S2 have 01 as last 2 inputs<br>S3 have $1^{\text {st }} 1$ of start of 100<br>S4 have 10 as last 2 inputs<br>S5 010 detected - 10 as last two inputs<br>S6 100 detected - output $\mathrm{Z} 2=1$<br>S7 after 100 - a 0 input<br>S8 after 100 - a 1 input<br>S9 after 100 - have 10 as last 2



## In S7 - have a 0

## - 0 input - Stay in S7

## - 1 input - transition to S8

| State | Meaning |
| :--- | :--- |
| S0 | Starting state |
| S1 | have $1^{\text {st }} 0$ of start of 010 |
| S2 | have 01 as last 2 inputs |
| S3 | have $1^{\text {st }} 1$ of start of 100 |
| S4 | have 10 as last 2 inputs |
| S5 | 010 detected -10 as last two inputs |
| S6 | 100 detected - output Z2 = 1 |
| S7 | after $100-$ a 0 input |
| S8 | after $100-$ a 1 input |
| S9 | after $100-$ have 10 as last 2 |



## In S8 - have xx01

## - 0 input - now have 10 - go to new state S9 <br> - 1 input - stay in S8

| State | Meaning |
| :--- | :--- |
| S0 | Starting state |
| S1 | have $1^{\text {st }} 0$ of start of 010 |
| S2 | have 01 as last 2 inputs |
| S3 | have $1^{\text {st }} 1$ of start of 100 |
| S4 | have 10 as last 2 inputs |
| S5 | 010 detected -10 as last two inputs |
| S6 | 100 detected - output Z2 = 1 |
| S7 | after $100-$ a 0 input |
| S8 | after $100-$ a 1 input |
| S9 | after $100-$ have 10 as last 2 |



## In state S9

## - 0 input - have seen 100 as last 3 - back to S6 <br> - 1 input - have $1^{\text {st }} 1$ of 100 - back to S8

| State | Meaning |
| :--- | :--- |
| S0 | Starting state |
| S1 | have $1^{\text {st }} 0$ of start of 010 |
| S2 | have 01 as last 2 inputs |
| S3 | have $1^{\text {st }} 1$ of start of 100 |
| S4 | have 10 as last 2 inputs |
| S5 | 010 detected -10 as last two inputs |
| S6 | 100 detected - output $\mathrm{ZZ}=1$ |
| S7 | after $100-$ a 0 input |
| S8 | after $100-$ a 1 input |
| S9 | after $100-$ have 10 as last 2 |



## Have seen contrast of Mealy/Moore

- Worked the development of a Mealy and Moore machine for the same specification
- Mealy - 8 states
- Moore - 10 states
- Machine has property that once certain conditions are met - a group of states can never be reached again. This type of machine is hard to test given the property of observeablilty.

